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**Proof-of-Concept (PoC) Report
Anti-Debugging and Anti-Emulation**

**For
SIRIUS Task Order PIQUE**

**Submitted to:
U.S. Government**

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07 August 2015

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(U) Executive Summary

(U) This report is an overview of techniques for detecting debuggers and emulation environments. This report represents a PoC delivery for July.

(U) Anti-Debugging

(U) There are a number of techniques to determine if a debugger is attached, including the use of Windows APIs, manually checking memory for debugger artifacts, and searching the system for forensics evidence of a debugger.

(U) Using Windows APIs

(U) The following Windows APIs can be used to determine if the application is being debugged:

IsDebuggerPresent()

CheckRemoteDebuggerPresent()

NtQueryInformationProcess()

CheckRemoteDebuggerPresent()

OutputDebugString()

(U) These function calls are easily implemented in code. However, like many APIs, they are easily hooked to provide a false answer as to whether the application is being debugged and therefore many malware authors prefer to manually check memory structures for the presence of a debugger, which is why we don't spend time on them in this report.

(U) Manually Checking Memory Structures

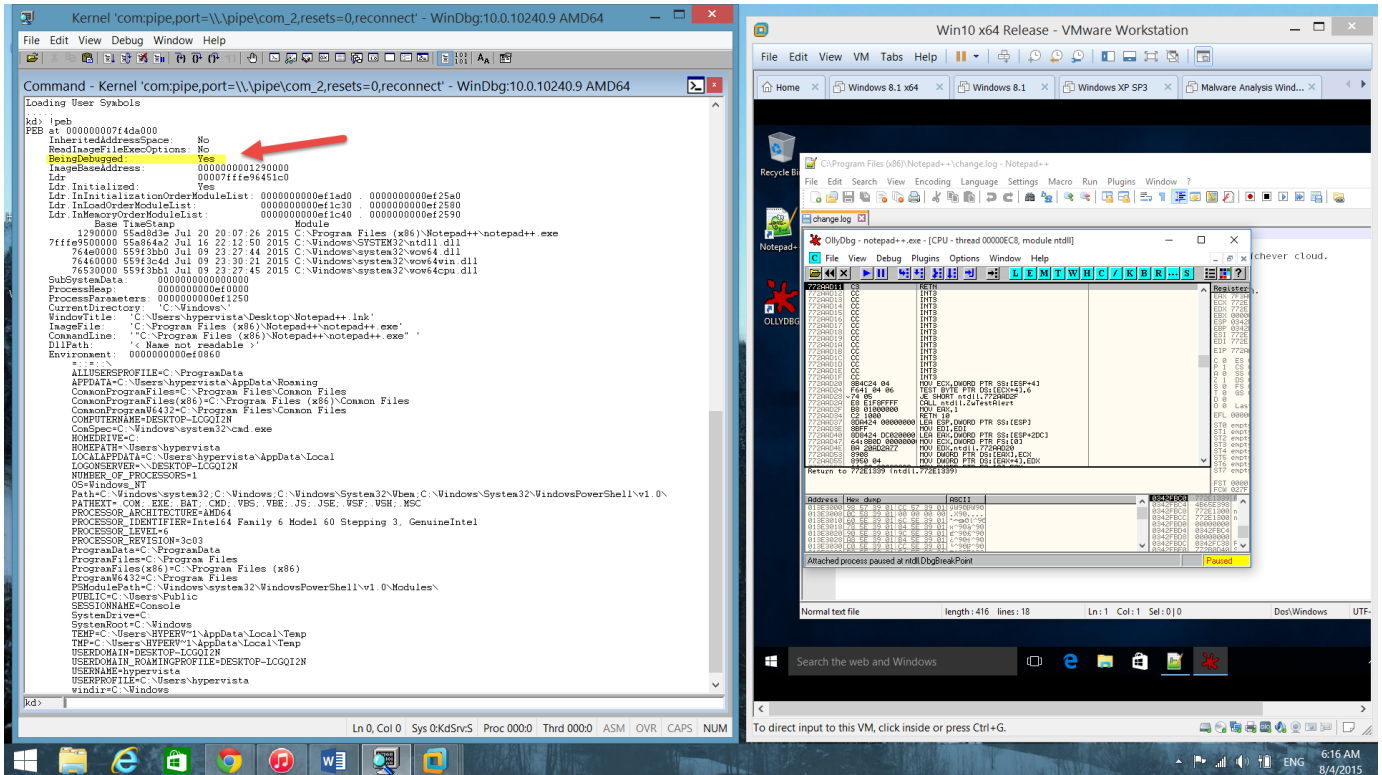
(U) Because Windows APIs can be hooked to return false information about whether or not the application is being debugged, it is sometimes preferable to manually check structures such as the PEB to determine if a debugger is present.

(U) Checking the PEB BeingDebugged Flag

(U) Windows maintains a Process Execution Block (PEB) structure for each process running. The PEB contains all user-mode parameters of the running process, including a flag relating to whether the process is being debugged or not (BeingDebugged). Figure 1 shows the Windows 10 x64 Enterprise PEB structure with the Notepad++ application being debugged using Ollydbg. The BeingDebugged flag can be checked programmatically by malware.

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(U) Figure 1. Windows 10 x64 Enterprise PE Structure - Application Debug

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(U) Program flow of a routine for checking the PEB BeingDebugged flag is represented in Figure 2:

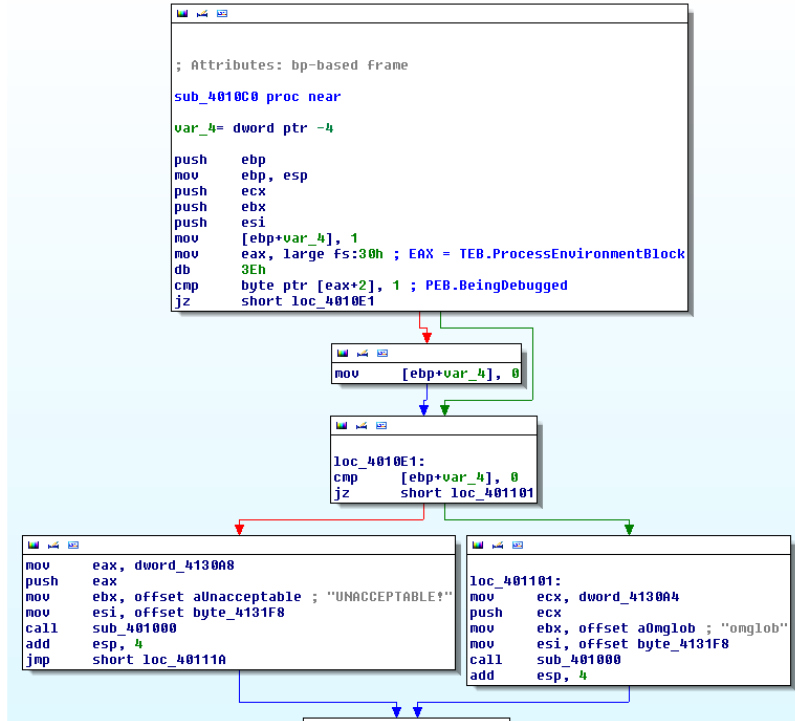


Figure 2. Programmatically Checking the BeingDebugged Flag

(U) Brief code explanation:

```
mov  eax, large fs:30h    ; move PEB structure (fs:30h) into EAX
cmp  byte ptr [eax+2], 1  ; check to see if the BeingDebugged flag is 1
iz   short loc_4010F1     ; if byte ptr [eax+2] = 1 it's being debugged and imn not made
```

(U) NOTE: there are several Ollydbg plug-in modules that change the BeingDebugged flag in order to ‘trick’ malware into believing it is not being debugged. The Ollydbg plug-ins that alter the BeingDebugged flag are: **Hide Debugger, Hidedebug, and PhantOm.**

(U) Checking the PEB ProcessHeap Flag

(U) Within the PEB is a data structure called ProcessHeap located at offset 0x18 in the PEB as shown in Figure 3. The first heap (at ef0000 in this example) contains a header with fields used to tell the kernel whether the heap was created within a debugger. The flags we are interested in are **ForceFlags** and **Flags**, which can be seen in Figure 4. The ForceFlags flag is a more reliable flag to check because Flags is usually either set to the value of ForceFlag or XORed with 2, as it is in our example as can be seen in Figure 4.

(U) The ForceFlags flag offset in the heap will differ depending on the OS version as shown in Table 1.

Table 1. ForceFlags Offset by Windows OS Version

OS Version	ForceFlags Offset
Windows XP 32-bit	0x10
Windows 7 32-bit	0x44
Windows 7 64-bit	0x74
Windows 8 32-bit	0x44
Windows 8 64-bit	0x74
Windows 10 32-bit	0x44
Windows 10 64-bit	0x74

(U) An example of an assembly language routine to manually check the ForceFlags flag is:

```

mov eax, large fs:30h           ; load PEB structure into EAX
mov eax, dword ptr [eax+18h]   ; Go to ProcessHeap offset in PEB
cmp dword ptr ds:[eax+74h], 0  ; Check to see if ForceFlags is 0 (offset varies per Table 1)
jne DebuggerDetected         ; DebuggerDetected routine defined elsewhere
    
```

```

kd> !peb
PEB at 000000007f4da000
  InheritedAddressSpace: No
  ReadImageFileExecOptions: No
  BeingDebugged: Yes
  ImageBaseAddress: 000000001290000
  Ldr: 00007ffe96451c0
  Ldr.Initialized: Yes
  Ldr.InInitializationOrderModuleList: 000000000ef1ad0 . 000000000ef25a0
  Ldr.InLoadOrderModuleList: 000000000ef1c30 . 000000000ef2580
  Ldr.InMemoryOrderModuleList: 000000000ef1c40 . 000000000ef2590
  Base TimeStamp: Module
    1290000 55ad8d3e Jul 20 20:07:26 2015 C:\Program Files (x86)\Notepad++\notepad++.exe
    7fffe9500000 55a864a2 Jul 16 22:12:50 2015 C:\Windows\SYSTEM32\ntdll.dll
    764e0000 559f3bb0 Jul 09 23:27:44 2015 C:\Windows\system32\wow64.dll
    76460000 559f3c4d Jul 09 23:30:21 2015 C:\Windows\system32\wow64win.dll
    76530000 559f3b01 Jul 09 23:27:45 2015 C:\Windows\system32\wow64cpu.dll
  SubSystemData: 0000000000000000
  ProcessHeap: 000000000ef0000 ←
  ProcessParameters: 000000000ef1250
  CurrentDirectory: C:\Windows\
  WindowTitle: C:\Users\hypervista\Desktop\Notepad++.lnk'
  ImageFile: C:\Program Files (x86)\Notepad++\notepad++.exe'
  CommandLine: "C:\Program Files (x86)\Notepad++\notepad++.exe"
  DllPath: %* Name not readable >
  Environment: 000000000ef0860
    =: =: \
    ALLUSERSPROFILE=C:\ProgramData
    APPDATA=C:\Users\hypervista\AppData\Roaming
    CommonProgramFiles=C:\Program Files\Common Files
    CommonProgramFiles(x86)=C:\Program Files (x86)\Common Files
    CommonProgramW6432=C:\Program Files\Common Files
    COMPUTERNAME=DESKTOP-ICGQI2N
    ComSpec=C:\Windows\system32\cmd.exe
    HOMEDRIVE=C:
    HOMEPATH=Users\hypervista
    LOCALAPPDATA=C:\Users\hypervista\AppData\Local
    LOGONSERVER=\\DESKTOP-ICGQI2N
    NUMBER_OF_PROCESSORS=1
    OS=Windows_NT
    Path=C:\Windows\system32;C:\Windows;C:\Windows\System32\Wbem;C:\Windows\System32\WindowsPowerShell\v1.0\
    PATHEXT=.COM;.EXE;.BAT;.VBS;.VBE;.JS;.JSE;.WSF;.WSH;.MSC
    PROCESSOR_ARCHITECTURE=AMD64
    PROCESSOR_IDENTIFIER=Intel64 Family 6 Model 60 Stepping 3, GenuineIntel
    PROCESSOR_LEVEL=6
    PROCESSOR_REVISION=3c03
    ProgramData=C:\ProgramData
    ProgramFiles=C:\Program Files
    ProgramFiles(x86)=C:\Program Files (x86)
    ProgramW6432=C:\Program Files
    PSModulePath=C:\Windows\system32\WindowsPowerShell\v1.0\Modules\
    PUBLIC=C:\Users\Public
    SESSIONNAME=Console
    SystemDrive=C:
    SystemRoot=C:\Windows
    TEMP=C:\Users\HYPERV*1\AppData\Local\Temp
    TMP=C:\Users\HYPERV*1\AppData\Local\Temp
    USERDOMAIN=DESKTOP-ICGQI2N
    USERDOMAIN_ROAMINGPROFILE=DESKTOP-ICGQI2N
    USERNAME=hypervista
    USERPROFILE=C:\Users\hypervista
    windir=C:\Windows
    
```

Figure 3. ProcessHeap in the PEB at offset 0x18


```

Command - Kernel 'compipe,port=\\.\pipe\com_2,reset=0,reconnect' - WinDbg:10.0.10240.9 AMD64
kd> !heap
Heap Address      NT/Segment Heap
ef0000            NT Heap
b80000            NT Heap
kd> dt _heap ef0000
ntdll!_HEAP
+0x00 Segment          : HEAP_SEGMENT
+0x00 Entry            : _HEAP_ENTRY
+0x00 Signature       : 0xfffffee
+0x00 Flags           : 2
+0x00 SegmentListEntry : LIST_ENTRY [ 0x00000000`00ef0120 - 0x00000000`00ef0120 ]
+0x00 Heap            : 0x00000000`00ef0000 _HEAP
+0x00 BaseAddress     : 0x00000000`00ef0000 Void
+0x00 NumberOfPages   : 0x1
+0x00 FirstEntry      : 0x00000000`00ef0710 _HEAP_ENTRY
+0x00 LastValidEntry  : 0x00000000`00ef0000 _HEAP_ENTRY
+0x00 NumberOfUnCommittedPages : 0xa
+0x00 NumberOfUnCommittedRanges : 1
+0x00 SegmentAllocatorBackTraceIndex : 0
+0x00 Reserved       : 0
+0x00 UCRSegmentList : LIST_ENTRY [ 0x00000000`00ef4fe0 - 0x00000000`00ef4fe0 ]
+0x00 Flags           : 2
+0x00 ForceFlags     : 0
+0x00 CompatibilityFlags : 0
+0x00 EncodeFlagMask : 0x100000
+0x00 Encoding        : HEAP_ENTRY
+0x00 Interceptor     : 0
+0x00 VirtualMemoryThreshold : 0xf00
+0x00 Signature       : 0xeeffeeff
+0x00 SegmentReserve  : 0x100000
+0x00 SegmentCommit  : 0x2000
+0x00 DeCommitFreeBlockThreshold : 0x100
+0x00 DeCommitTotalFreeThreshold : 0x1000
+0x00 TotalFreeSize   : 0x27
+0x00 MaximumAllocationSize : 0x00007fff`ffffefff
+0x00 ProcessHeapsListIndex : 1
+0x00 HeaderValidateLength : 0x298
+0x00 HeaderValidateCopy : (null)
+0x00 NextAvailableTagIndex : 0
+0x00 MaximumTagIndex : 0
+0x00 TagEntries      : (null)
+0x00 UCRList         : LIST_ENTRY [ 0x00000000`00ef4fd0 - 0x00000000`00ef4fd0 ]
+0x00 AlignRound      : 0x1f
+0x00 AlignMask       : 0xfffffff0
+0x00 VirtualAllocdBlocks : LIST_ENTRY [ 0x00000000`00ef0110 - 0x00000000`00ef0110 ]
+0x00 SegmentList     : LIST_ENTRY [ 0x00000000`00ef0018 - 0x00000000`00ef0018 ]
+0x00 AllocatorBackTraceIndex : 0
+0x00 NonDedicatedListLength : 0
+0x00 BlocksIndex     : 0x00000000`00ef02c0 Void
+0x00 UCRIndex        : (null)
+0x00 PseudoTagEntries : (null)
+0x00 FreeLists       : LIST_ENTRY [ 0x00000000`00ef26a0 - 0x00000000`00ef26a0 ]
+0x00 LockVariable    : 0x00000000`00ef0298 _HEAP_LOCK
+0x00 CommitRoutine   : 0x5f3ab8db`818791ff long +5f3ab8db818791ff
+0x00 FrontEndHeap    : (null)
+0x00 FrontHeapLockCount : 0
+0x00 FrontEndHeapType : 0
+0x00 RequestedFrontEndHeapType : 0
+0x00 FrontEndHeapUsageData : 0x00000000`00ef0720 -> 0
+0x00 FrontEndHeapMaximumIndex : 0x00
+0x00 FrontEndHeapStatusBitmap : [129] **
+0x00 Counters        : _HEAP_COUNTERS
+0x00 TuningParameters : _HEAP_TUNING_PARAMETERS
kd>

```

Figure 4. ForceFlags Flag at offset 0x74 (Windows 10 64-bit)

(U) Checking the PEB NTGlobalFlag Value

(U) Another PEB element that contains information about whether the application was launched under a debugger is the NTGlobalFlag. The NTGlobalFlag element offset in the PEB differs between the 32-bit and 64-bit versions of Windows, as shown in Table 2.

Table 2. NTGlobalFlag Offset in PEB by OS Variety

Windows Variety	NTGlobalFlag Offset in PEB
32-bit Versions (Win 8, 8.1, and 10)	0x068
64-bit Versions (Win 8, 8.1, and 10)	0x0bc

(U) If the value of the NTGlobalFlag is 0x70, the application is being debugged as can be seen in Figure 5. The value of 0x70 is a combination of three separate flags being set when a heap is created by a debugger. The three flags set when a heap is created by a debugger are FLG_HEAP_ENABLE_TAIL_CHECK (0x10), FLG_HEAP_ENABLE_FREE_CHECK (0x20), and FLG_HEAP_VALIDATE_PARAMETERS (0x40).

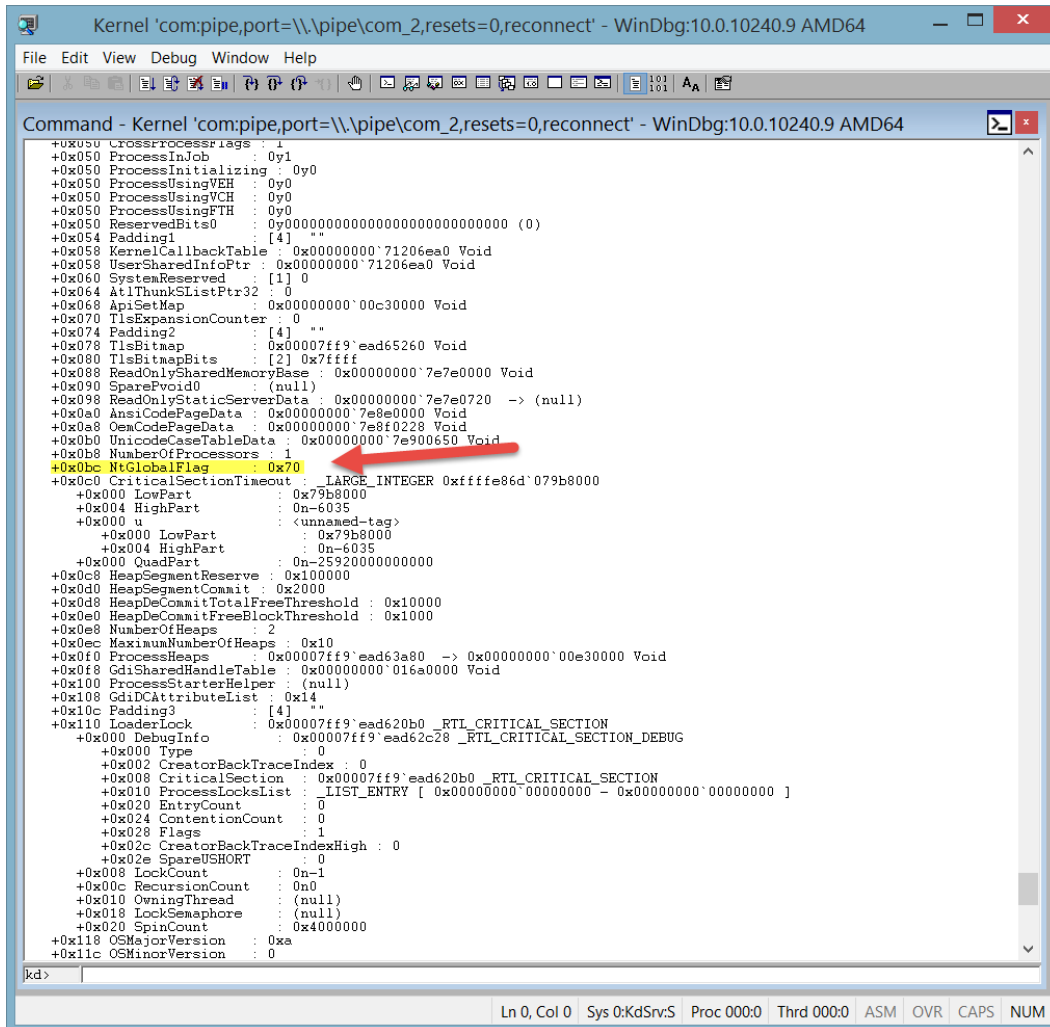


Figure 5. NTGlobalFlag - 64-bit Application Being Debugged

(U) An assembly language example of manually checking the value of the NTGlobalFlag is:

```
mov eax, large fs:30h           ; move PEB structure into EAX
cmp dword ptr ds:[eax+bch], 70h ; check to see if NTGlobalFlag is 0x70
jz DebuggerDetected           ; DebuggerDetected function defined elsewhere
```

(U) Identifying Debugger Behavior

(U) Because some debugging activities by necessity modify the code at debugger runtime, these alterations of the code can be detected. Debugger generated code modifications include insertion of INT instructions (not just INT 3), these INT instructions can be scanned for. Malware can use checksums on their code to determine if the running code has been altered in any way, presumably by a debugger. Lastly, malware can also perform timing checks because processes run slower when being debugged.

(U) Checking to See if SeDebugPrivilege is Set

(U) By default, a process has SeDebugPrivilege disabled. When the process is loaded by a debugger, SeDebugPrivilege is enabled. Malware will check to see if SeDebugPrivilege has been enabled by trying to open the CSRSS.EXE process and if it is able to open it the process is running under a debugger.

(U) Scanning for INT

(U) INT 3 is a software interrupt that debuggers insert into running code, replacing an existing instruction, which calls the debug exception handler, i.e. sets a breakpoint. The opcode for INT 3 is 0xcc. In addition to inserting an INT 3 opcode into running code, debuggers also insert INT <immediate value> in some cases. The opcode for INT

```
mov ecx, 400h           ; loop counter
mov eax, 0cch           ; INT 3 (0cch) search value
repne scasb            ; string search command (look for INT 3 (0cch))
jz DebuggerDetected    ; DebuggerDetected function is defined elsewhere
```

<immediate value> is 0xcd. A common method malware uses to detect if it's running under the control of a debugger is to scan for these opcodes. A rough assembly language routine to scan for INT 3 is:

(U) Code Checksums

(U) Some malware samples calculates checksums of specific sections of its code, either CRC or MD5 to detect debugger modification of the code to implement breakpoints. This technique is less common than INT scanning, but just as effective.

(U) Timing Checks

(U) Because processes run substantially slower under a debugger (think single-stepping through code), timing checks is a very effective and popular way malware authors check for the presence of a debugger. There are a few methods for conducting timing checks for the presence of a debugger:

Take a timestamp, perform some specific operations and take another timestamp, calculate the time difference and make a judgment about whether or not the time difference is outside the bounds of normal time to conduct the operations.

Take a timestamp before and after causing an exception, calculate the time difference and make a judgment about whether the time to respond to the exception is well outside the normal time require to handle an exception.

(U) The use of the rdtsc command is quite popular among malware authors as a time checking function. In addition to rdtsc, the Windows APIs QueryPerformanceCounter() and GetTickCount are also heavily used as well.

(U) Checking the Number of Kernel DebugObjects

(U) When an application is being debugged, a DebugObject is created in the kernel. The number and type of Objects created can be retrieved via the NtQueryObject() API, which returns an OBJECT_ALL_INFORMATION structure. The OBJECT_ALL_INFORMATION structure is searched for the string, "DebugObject" and is checked for a non-zero value, indicating the presence of a debugger.

(U) Checking for a Debugger Window

(U) Most debuggers create windows, which can be detected. For example, Windbg creates WinDbgFrameClass and Ollydbg creates OLLYDBG. The Windows APIs FindWindow() and FindWindowEx() can be used to search for the debugger windows, indicating the presence of a debugger.

(U) Providing an Invalid ASCII String to OutputDebugStringA

(U) Calling OutputDebugStringA() with an invalid ASCII string will normally return a value of 1. If the process is being run under the control of a debugger, the return value when providing an invalid ASCII string to OutputDebugStringA() is the address of the string passed in as a parameter, indicating the presence of a debugger.

(U) Using the Stack Segment Register and Checking Trapflag

(U) This is an anti-tracing technique. If a debugger is tracing over a sequence of instructions that includes pop ss and pushf instructions, the debugger will not be able to unset the trapflag in the pushed value on the stack. The protection checks for the trapflag and if set it indicates the presence of a debugger. For example:

```
push  ss
pop   ss
pushf
```

(U) When tracing over the pop ss instruction in this example code, the next instruction will be executed but the debugger will not break on it, therefore stopping on the following instruction, nop in this case. This rare anti-debugging trick

```
push  ss
; junk code here
pop   ss
pushf
; junk code here
pop   eax
and   eax,100h
or    eax,eax
```

has been seen in the wild as follows:

(U) The trick here is that if the debugger is tracing over this sequence of instructions, popf will be executed implicitly and the debugger will not be able to unset the trapflag. The malware checks for the trapflag and debugger is present if found.

(U) Trolling the Debugger

(U) In addition to detecting the presence of a debugger and responding accordingly, by either exiting without installing malicious code or by presenting benign behavior, some malware has been observed interfering with debugger functionality to make a malware analyst's job more difficult. While we're pretty certain the Sponsor would not use such tactics, we're including these techniques for completeness.

(U) Modifying the SEH Chain

(U) Modifying the SEH chain can be used as an anti-disassembly technique as well as an anti-debugging technique. Exception-based detection relies on the fact that debuggers will trap the exception and not immediately pass it to the process for handling. If the debugger fails to pass the exception to the process for handling, as most will do, that can be detected within the exception-handling mechanism and a determination that a debugger is running can be made.

(U) Inserting INT Commands

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(U) This technique is more an obnoxious annoyance than anything. Some malware samples have been observed inserting INT 3 commands into code to cause the debugger to break, or in some cases stop all together.

(U) Inserting an INT 1 will invoke single-step mode inside the debugger.

(U) Some malware specimen insert INT 2D commands to cause the kernel debugger to break, similar to the insertion of INT 3 into application code.

(U) Following is an assembly language example of malware inserting an INT 3 command to check for the presence of a debugger. This example sets EAX to 0xffffffff inside the exception handler to signify the exception handler had been called. If EAX is not 0xffffffff after the INT 3 call, then a debugger is present.

```
;set exception handler
    push    .exception_handler
    push    dword [fs:0]
    mov     [fs:0], esp

;reset flag (EAX) invoke int3
    xor     eax, eax
    int3

;restore exception handler
    pop     dword [fs:0]
    add     esp, 4

;check to see if the flag has been set
    test    eax, eax
    je     .debugger_found

;exception_handler
```

(U) Inserting In-Circuit Emulator

(ICE) Breakpoints

(U) Inserting an ICE breakpoint, iceb (opcode 0xf1) generates a single-step exception and the debugger will think it's a normal exception and not execute the established exception handler. Malware can take advantage of this fact by using the exception handler for its normal execution flow, which would be interrupted if a debugger is attached thereby hiding the malicious code from examination.

(U) Anti-Emulation

(U) Virtual environments, virtual machines and commercial sandbox technologies built upon them are collectively known as emulation environments. Emulation environments pose a risk to malware because they are used by system defenders to analyze malware at runtime. Malware authors have spent a lot of resources and written many routines to detect emulation environments in an attempt to protect their malware. However, in recent years we've seen a decrease in anti-emulation techniques in malware and attribute this trend to the fact that emulation environments have become so pervasive (think Cloud) that the presence of an emulation environment no longer automatically means it's a malware analysis platform or that it's not a valid target.

(U) Detecting VMWare Artifacts

(U) VMWare leaves many detectable artifacts when installed on a system. Malware can use these artifacts left in the file system, registry, and process listing. The following were run on a Windows 10 Virtual Machine image.

(U) Using *net start | findstr VMWare*

(U) A quick and easy way to determine if the VMWare Tools Service is running on the system is to open a command window and type *net start | findstr VMWare* as shown in Figure 6.

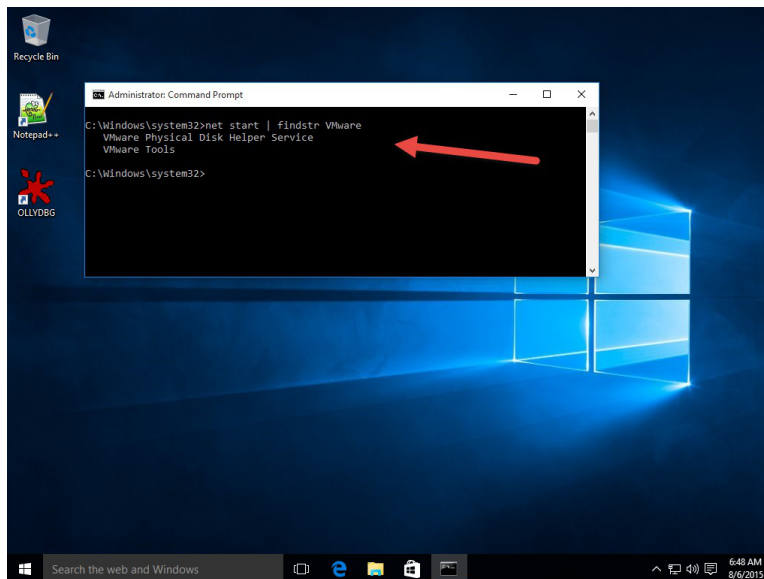


Figure 6. Using cmd.exe to Find the VMWare Service

(U) Searching the File System

(U) Another quick and easy way to determine if VMWare has been installed on the target is to look in the file system (./Program Files/VMWare) as shown in Figure 7. This file system search was conducted on a Windows 10 VMWare machine, not the host machine.

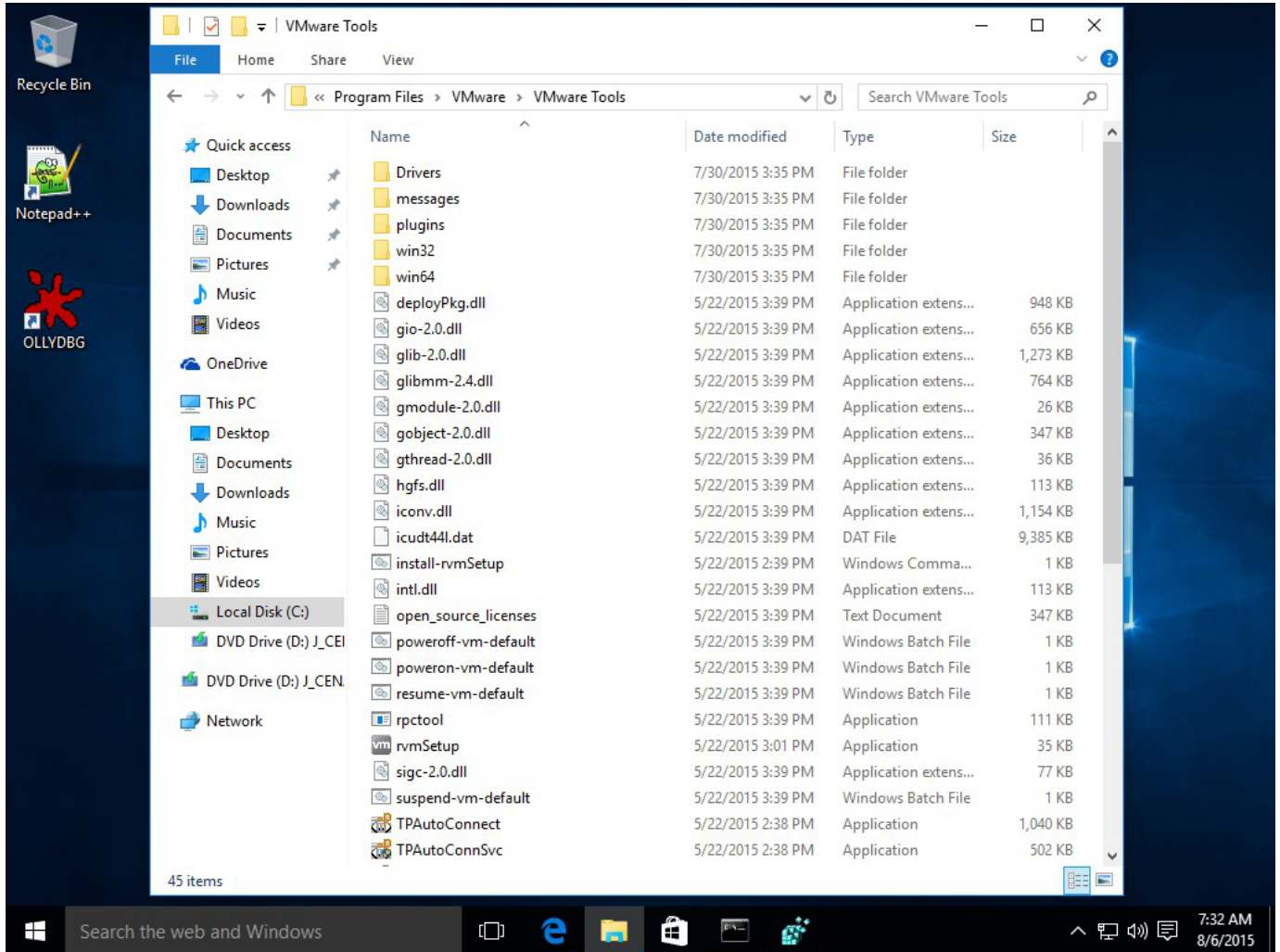


Figure 7. Searching the File System for VMware

(U) Search the Registry for 'VMWare'

(U) Searching the registry for the string 'VMWare' is another simple and easy way to look for VMWare installation artifacts on the target machine. The following registry search, as shown in Figure 8, was conducted on a Windows 10 guest Virtual Machine, not the host machine.

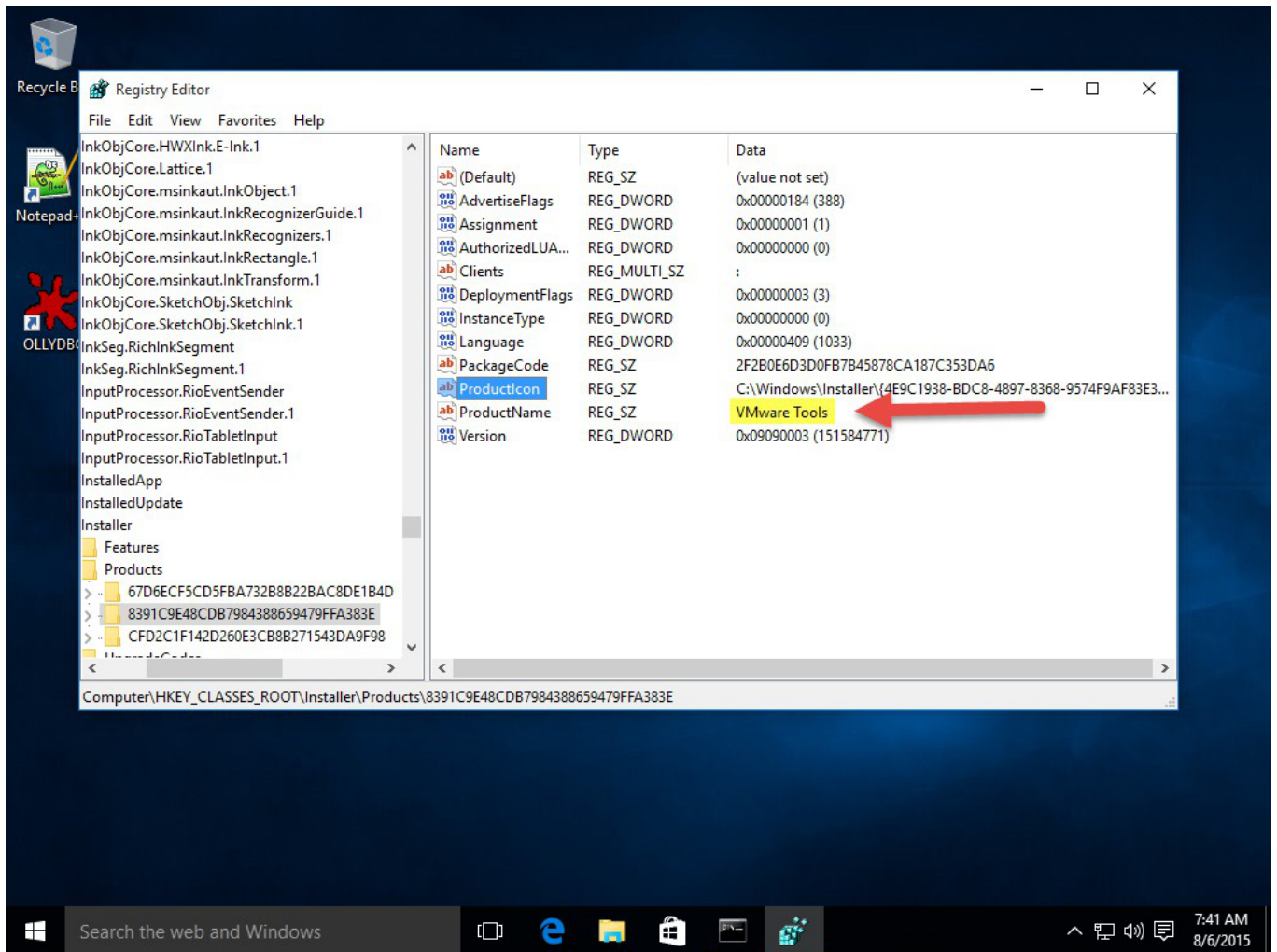


Figure 8. Searching the Registry for 'VMWare'

(U) Checking the MAC for Leading 00:0C:29

(U) The first three groups in a MAC address identify the manufacturer of the network device.

VMWare's default MAC address identifier groups are 00:0C:29. An easy check is to open a command prompt and type *getmac* and look for the VMWare MAC identifier as shown in Figure 9.

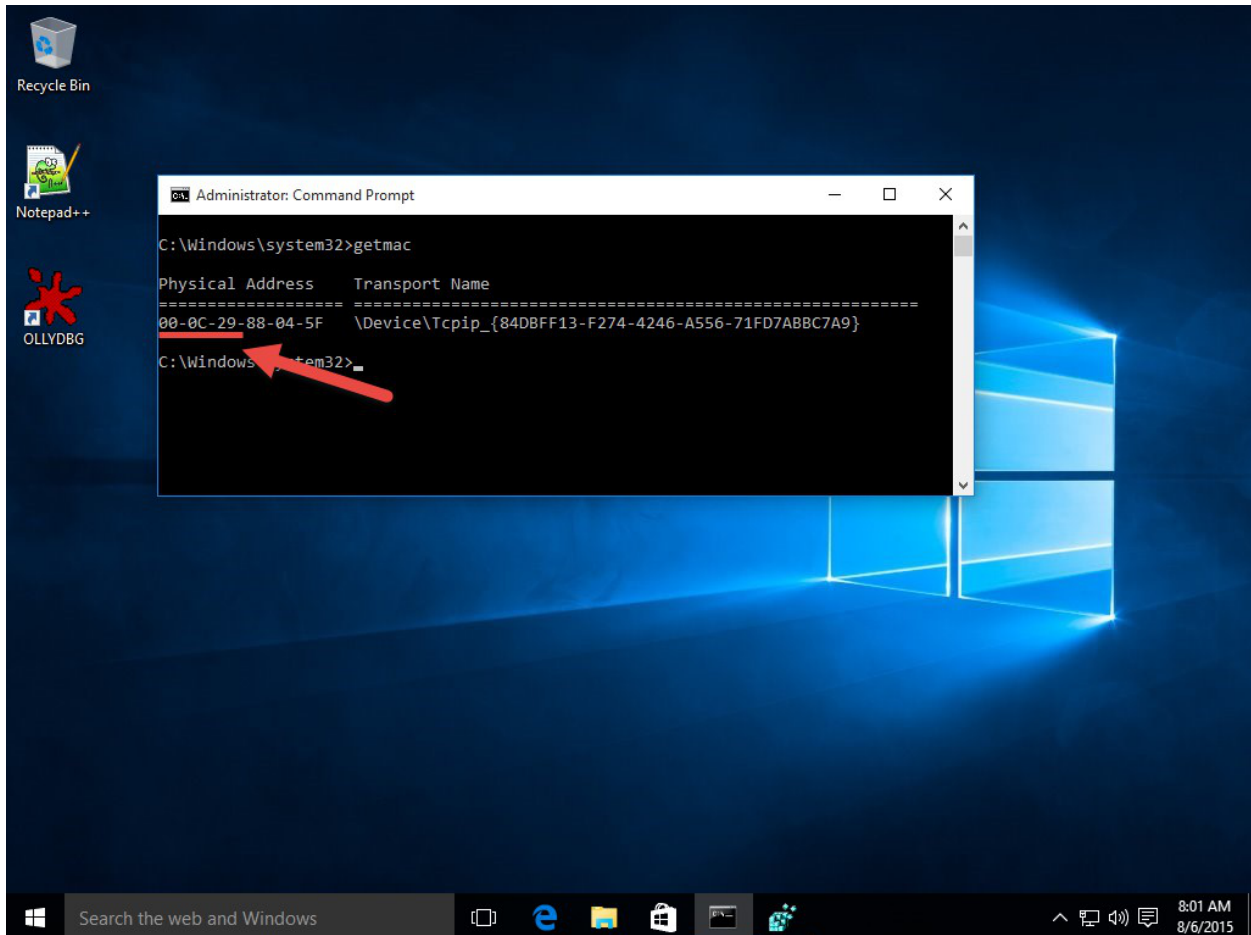


Figure 9. VMWare MAC Identifier 00-0C-29

(U) Using Sensitive Instructions to Detect VMWare

(U) Because not all instructions can be virtualized, VMWare uses what is known as binary translation for these ‘problematic’ instruction. Binary translation traps the problematic instructions, and essentially changes context to handle the instruction on the host processor, returning the result and resuming the virtualization. Naturally, this causes a significant performance hit and timing methods can be used to determine if the target is running in VMWare.

(U) Some instructions return different results if run in VMWare as compared to running on native hardware. These instructions can be used to determine if a VMWare machine is present.

(U) The Red Pill Anti-VM Technique

(U) The Red Pill anti-VM technique uses the sidt instruction to get the value of the IDTR register. The VM has to relocate the guest IDTR to avoid conflict with the host’s IDTR. Because the VM is not notified when the VM runs the sidt instruction, the IDTR for the VM is returned.

The fifth byte of the IDTR returned contains the start of the base memory address. VMWare returns 0xFF in the fifth byte. The Red Pill tests for this discrepancy to detect VMWare. It should be noted that this technique only works reliably on single processor machines (VMWare default is single processor).

(U) The No Pill Anti-VM Technique

(U) The No Pill anti-VM technique uses the sgdt and sltd and relies on the fact the Local Descriptor Table (LDT) structure is assigned to a processor, not the OS. Because Windows does not use the LDT structure but VMWare provides support for it the table will vary in predictable ways. The location of LDT on the host machine will be zero while the location of the LDT of a VMWare guest machine will be a non-zero value. A simple check for a non-zero return from either sgdt or sltd will indicate the OS is a VM image.

(U) Checking the I/O Communications Port

(U) VMWare uses a virtual I/O port with a specific ‘magic number’ for communications between the host machine and the guest machine. The port can be queried and compared with the magic number to identify VMWare. The magic number, in hex is 0x564D5868, converts to ASCII “VMXh”.

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Pique Proof-of-Concept (PoC) Report Anti-Debugging and Anti-Emulation

(U) This technique relies on the x86 instruction *in*, which copies data from the I/O port specified by the source operand to a memory address specified by the destination operand. VMWare monitors the use of the *in* operator and captures the I/O destination for the communication port 0x5668 (VX). Therefore, the second operand needs to be loaded with VX in order to check for VMWare, which only happens when EAX holds the magic number 0x564D5868. ECX must contain a value corresponding to the action you want to perform on the port. The value 0xA means, “get the VMWare version type”, and 0x14 means, “get the memory size.” Either can be used to detect VMWare.

Timing-based VM Detection

(U) As mentioned earlier, virtualization environments cannot emulate every instruction. Some instructions are trapped in the kernel, the virtualization environment is halted, the instruction is handled by the host processor and the result passed back to the emulation environment and the emulation environment is restarted. Naturally, all this causes quite a performance hit and that degradation of performance can be measured via timestamps.

(U) One of the problematic instructions is CPUID. The timing-based VM detection technique involves taking a timestamp, looping a large number of sequential CPUID calls (> 4000), taking another timestamp and calculating the difference. If running in an emulation environment, the time difference will be orders of magnitude greater than the time difference when running on a native host system.

(U) Resources

www.symantec.com/connect/articles/windows-anti-debug-reference

Andrew Honig and Michael Sikorski, Practical Malware Analysis, No Starch Press 2012

<http://www.aldeid.com/wiki/PEB-Process-Environment-Block/BeingDebugged>

<https://www.blackhat.com/presentations/bh-usa-07/Yason/Whitepaper/bh-usa-07-yason-WP.pdf>

<https://www.exploit-db.com/docs/34591.pdf>